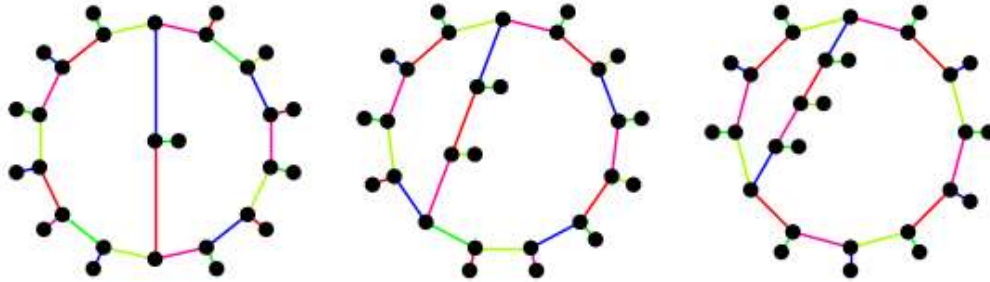


# Strong Chromatic Index

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## Problem and Results Summary

Let  $G$  be a graph. A  $k$ -strong-edge-coloring is a function  $c : E(G) \rightarrow \{1, \dots, k\}$  where every color-class forms an *induced matching*. That is, two edges of the same color are not incident and do not have adjacent endpoints. Equivalently, a  $k$ -strong-edge-coloring is a  $k$ -coloring of the square of the line graph. The *strong chromatic index*, denoted  $\chi_s(G)$ , is the minimum  $k$  such that a  $k$ -strong-edge-coloring exists for  $G$ .

Many people have studied the strong chromatic index of planar graphs which are subcubic (maximum degree at most three).

**Theorem.** (Kostochka et al. 2015+). If a graph  $G$  is planar and subcubic then  $\chi_s(G) \leq 9$ .

Other sparsity conditions can imply stronger bounds on the strong chromatic index. For instance, if we force the girth of a planar graph to be high, then it becomes easier to color.

**Theorem.** (Hocquard et al. [8]). Let  $G$  be a subcubic planar graph with girth  $g$ .

1. If  $g \geq 14$ , then  $\chi_s(G) \leq 6$ .
2. If  $g \geq 10$ , then  $\chi_s(G) \leq 7$ .
3. If  $g \geq 8$ , then  $\chi_s(G) \leq 8$ .

The potentially optimal number of colors for a graph with maximum degree  $\Delta$  is  $2\Delta - 1$ . In this scenario, there is a sequence of results.

**Theorem.** Let  $G$  be a graph with  $\Delta(G) \leq d$ .

1. (Borodin and Ivanova) If  $d = 3$ ,  $G$  is planar, and  $G$  has girth at least 41, then  $\chi_s(G) \leq 5$ .
2. (Borodin and Ivanova; implicit) If  $d = 3$ ,  $G$  has girth at least 9, and  $\text{mad}(G) < 2 + \frac{2}{23}$ , then  $\chi_s(G) \leq 5$ .
3. (Borodin and Ivanova) If  $d > 3$ ,  $G$  is planar, and  $G$  has girth at least  $40\lfloor \frac{d}{2} \rfloor + 1$ , then  $\chi_s(G) \leq 2d - 1$ .
4. (Borodin and Ivanova; implicit) If  $d > 3$ ,  $G$  has girth at least  $8\lfloor \frac{d}{2} \rfloor + 1$ , and  $\text{mad}(G) < 2 + \frac{2}{24\lfloor \frac{d}{2} \rfloor - 1}$ , then  $\chi_s(G) \leq 2d - 1$ .
5. (Chang, Montassier, Pêcher, Raspaud) If  $d > 3$ ,  $G$  is planar, and  $G$  has girth at least  $10d + 46$ , then  $\chi_s(G) \leq 2d - 1$ .
- 6.
7. (Chang, Montassier, Pêcher, Raspaud; implicit) If  $d > 3$ ,  $G$  has girth at least  $2d + 10$ , and  $\text{mad}(G) < 2 + \frac{2}{6d + 26}$ , then  $\chi_s(G) \leq 2d - 1$ .
8. (Wang and Zhao) If  $d > 3$ ,  $G$  is planar, and  $G$  has girth at least  $10d - 4$ , then  $\chi_s(G) \leq 2d - 1$ .
9. (Wang and Zhao) If  $d > 3$ ,  $G$  has girth at least  $2d - 1$ , and  $\text{mad}(G) < 2 + \frac{1}{3d - 2}$ , then  $\chi_s(G) \leq 2d - 1$ .
10. (DeOrsey, Diemunsch, Ferrara, Graber, Harke, Jahanebkam, Lidický, Nelsen, Stolee, Sullivan) If  $d = 3$ ,  $G$  is planar, and  $G$  has girth at least 30, then  $\chi_s(G) \leq 5$ .

11. (DeOrsey, Diemunsch, Ferrara, Graber, Harke, Jahanebkam, Lidický, Nelsen, Stolee, Sullivan)  
If  $d = 3$ ,  $G$  has girth at least 9, and  $\operatorname{mad}(G) < 2 + \frac{1}{7}$ , then  $\chi_s(G) \leq 5$ .
12. (DeOrsey, Diemunsch, Ferrara, Graber, Harke, Jahanebkam, Lidický, Nelsen, Stolee, Sullivan)  
If  $d = 4$ ,  $G$  is planar, and  $G$  has girth at least 28, then  $\chi_s(G) \leq 7$ .
13. (DeOrsey, Diemunsch, Ferrara, Graber, Harke, Jahanebkam, Lidický, Nelsen, Stolee, Sullivan)  
If  $d = 4$ ,  $G$  has girth at least 7, and  $\operatorname{mad}(G) < 2 + \frac{2}{13}$ , then  $\chi_s(G) \leq 7$ .

The theorems above are implied by a sparsity condition using the maximum average degree of a graph (and some forbidden subgraphs). A related result implies that the list-coloring version of strong chromatic index is bounded by five for subcubic planar graphs with girth at least 41. All results use discharging in some manner, either explicitly or implicitly.

## Research Papers

Philip DeOrsey, Jennifer Diemunsch, [Michael Ferrara](#), Nathan Graber, Stephen G. Harke, Sogol Jahanebkam, [Bernard Lidický](#), Luke Nelsen, Derrick Stolee, Eric Sullivan, [On the Strong Chromatic Index of Sparse Graphs](#)  
[Web Version](#)

## Software

There are three main software components.

- [SCI\\_magma.txt](#) is a list of [Magma](#) commands to discover the coefficient of a polynomial to use in a Combinatorial Nullstellensatz argument. (Written by Phil DeOrsey)
- [SCI\\_Sage\\_code.sage](#) is a list of [Sage](#) commands to discover the coefficient of a polynomial to use in a Combinatorial Nullstellensatz argument. (Written by Stephen G. Hartke)
- [StrongEdgeColorings.sws](#) is a [Sage](#) worksheet that contains many helpful routines, including constructing the configurations. It also tests the base cases of the main theorem regarding strong 5-edge-colorings.
- [strong\\_edge\\_reducible.c](#) is a C program that tests if a configuration is reducible.

To compile, place this file in the `nauty/` folder after unpacking the `nauty` source code, available at <http://cs.anu.edu.au/~bdm/nauty/>

Compile using the following command:

```
gcc -o strong_edge_reducible strong_edge_reducible.c nauty.o nauparse.o nautil.o gtools.o
```

Then, run : `strong_edge_reducible.exe < configurations.txt`

or (for 7 colors) : `strong_edge_reducible.exe -k 7 < configurations.txt`

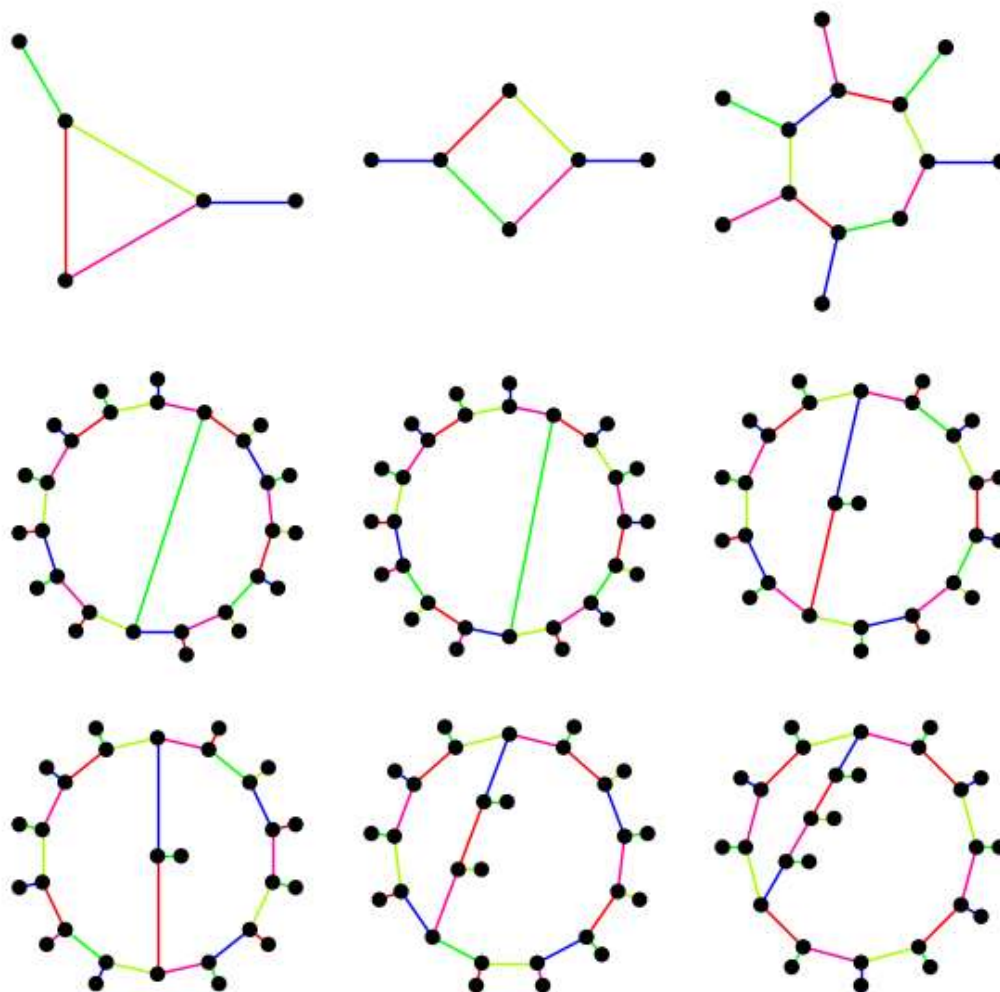
- [YColoring.sws](#) is a [Sage](#) worksheet that determines that no 3-vertex  $v$  can have  $N_3(v)$  consist entirely of vertices  $u$  with  $|\operatorname{Resp}(u)| \geq 12$ . (Written by Bernard Lidický)

## Data

- [configurations.txt](#) is the input file for `strong_edge_reducible.exe` containing all configurations, reducible or not.
- [ser\\_output.txt](#) is the output file reporting which configurations are reducible or not.
- [caterpillars-d4.txt](#) is the input file for `strong_edge_reducible.exe` containing a list of  $(t,4)$ -caterpillars.
- [caterpillars-d5.txt](#) is the input file for `strong_edge_reducible.exe` containing a list of  $(t,5)$ -caterpillars.
- [caterpillars-d6.txt](#) is the input file for `strong_edge_reducible.exe` containing a list of  $(t,6)$ -caterpillars.
- [reducible-d4y.txt](#) is the input file for `strong_edge_reducible.exe` containing a list of  $Y_4(t_1, t_2, t_3)$  configurations.
- [reducible-d4y-out.txt](#) is the output file for `strong_edge_reducible.exe -k 7 < reducible-d4y.txt` demonstrating which  $Y_4(t_1, t_2, t_3)$  configurations are 7-reducible.

## Gallery

Here are some strong edge colorings that were used as base cases in the proof.



## Support

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## Related Work

H. Hocquard, M. Montassier, A. Raspaud, and P. Valicov, On strong edge-colouring of subcubic graphs, *Discrete Applied Mathematics* 161 (2013) 2467-2479.

A.V. Kostochka, X. Li, W. Ruksasakchai, M. Santana, T. Wang, and G. Yu, Strong chromatic index of subcubic planar multigraphs, *manuscript in preparation*.